Principles of Software Construction: Objects, Design, and Concurrency

{Static & Dynamic} x {Typing & Analysis}

Claire Le Goues

Vincent Hellendoorn



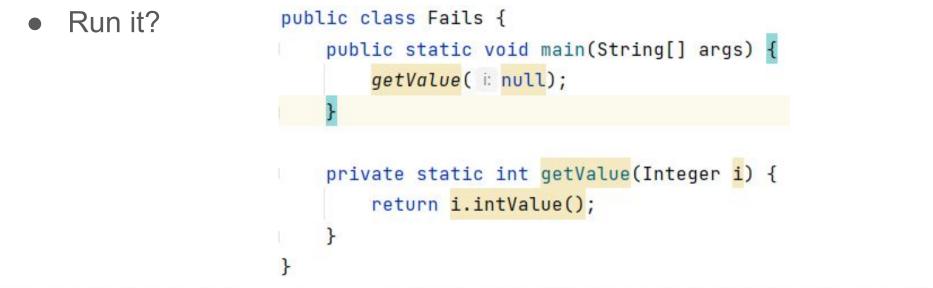


Is this code buggy?

private static int getValue(Integer i) { return i.intValue(); }



How Do You Find Bugs?



Exception in thread "main" java.lang.NullPointerException Create breakpoint : Cannot invoke "java.lang.Integer.intValue()" because "i" is null
 at misc.Fails.getValue(Fails.java:9)
 at misc.Fails.main(Fails.java:5)



17-214/514

How Else Can You Find Bugs?

```
public class Fails {
    public static void main(String[] args) {
        getValue( i: null);
    }
    private static int getValue(Integer i) {
        return i.intValue();
    }
}
```



}

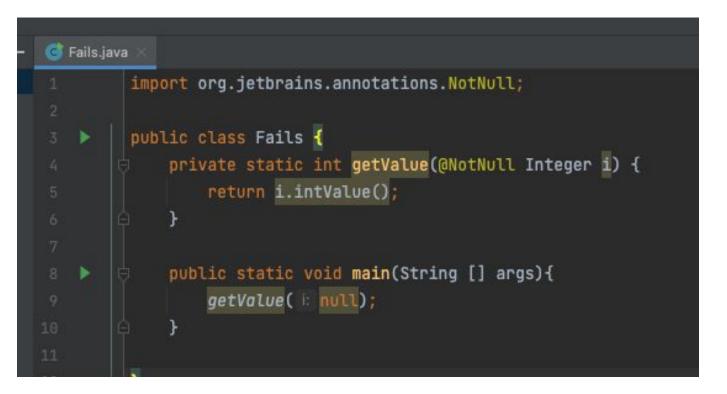


IntelliJ can look at this code and say:

```
public static void main(String[] args) {
    getValue( i: null);
}
Passing 'null' argument to parameter annotated as @NotNull
private static int getValue(Integer i) {
    return i.intValue();
}
```



(with annotations explicit)







How?

```
public static void main(String[] args) {
    getValue( i: null);
}
Passing 'null' argument to parameter annotated as @NotNull
private static int getValue(Integer i) {
    return i.intValue();
}
```





How?

- We know at *compile time* where getValue gets routed to
- getValue calls a method on i
- i can be null

```
public static void main(String[] args) {
    getValue( i: null);
}
Passing 'null' argument to parameter annotated as @NotNull
private static int getValue(Integer i) {
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```



What about JS?

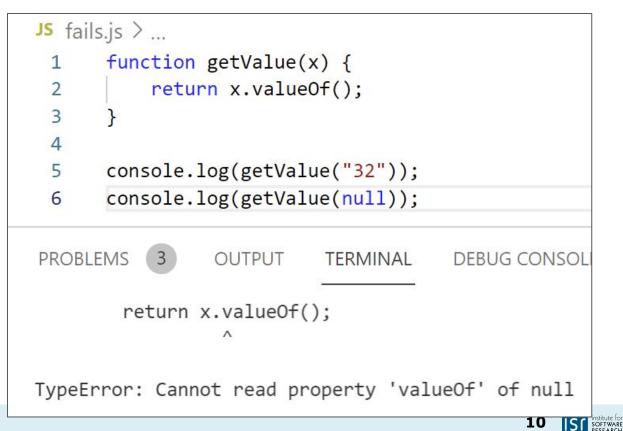
fails.js
function getValue(x) {
 return x.valueOf();
}





What about JS?

Run it: 🗸



```
Why no warning?
```

```
function getValue(x) {
    return x.valueOf();
}
```

console.log(getValue("32")); console.log(getValue(null));





Another Java vs JS Example

```
class Foo {
     constructor(x) {
          this.x = x;
 function bar(foo) {
     return foo.x;
 var foo = new Foo(3);
 console.log(bar(foo));
<sup>17</sup>console.log(bar(3)):
```

```
class Foo {
    int x;
    Foo(int x) {
        this.x = x;
public static void main(String[] args) {
    Foo foo = new Foo(\times 3);
    bar(foo);
    bar( foo: 3);
private static void bar(Foo foo) {
    System.out.println(foo.x);
l
```

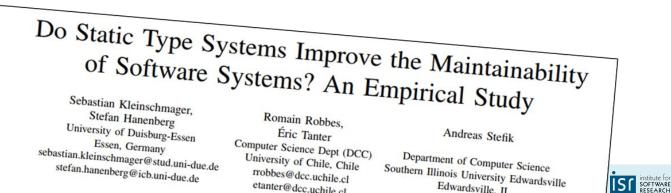
- The more knowledge we inject in the code, the more bugs we can catch at compile time
 - Types, nullity annotations, invariants
- At compile-time:
 - Dynamically typed languages assume nothing
 - Types exist only for values
 - Static typing is not completely precise either
 - Objects have declared types and run-time types
 - Different "strength" type systems



- The more knowledge we inject in the code, the more bugs we can catch at compile time
 - Types, nullity annotations, invariants
- Is it worth it?
 - Dynamic typing can severely limit inference
 - But... static types are a lot of work



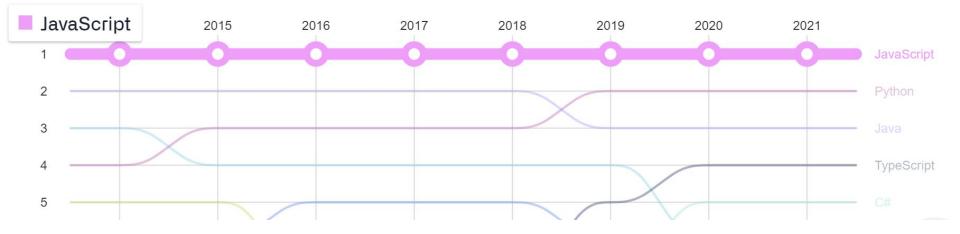
- The more knowledge we inject in the code, the more bugs we can catch at compile time
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Okay, but:

Top languages over the years

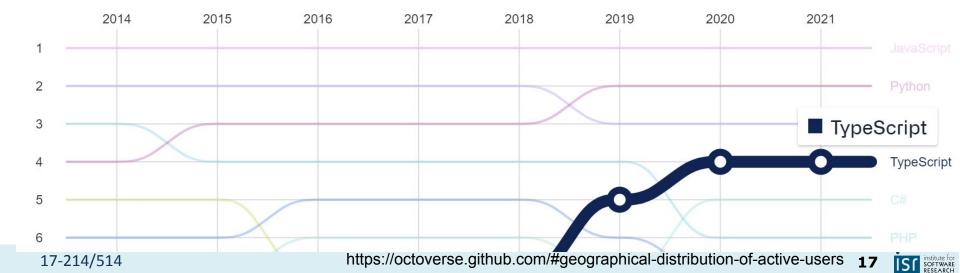




False Dichotomy?

Yes, but:

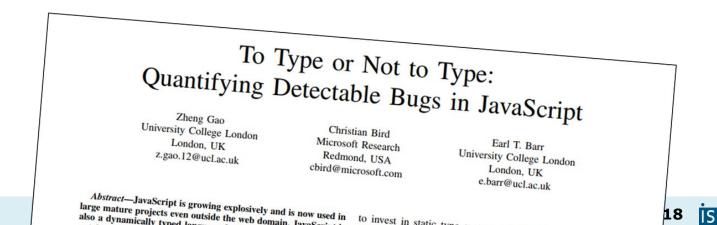
Top languages over the years



Partial Types

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- Low effort, some utility
 - Static types exist and are checked at compile-time
 - Dynamic types are used at run-time
 - So annotations get ignored!
 - Type checker can be shallow or deep; TS is shallow



SOFTWARE RESEARCH

Types in TypeScript

function getValue(x: number) {
 return x.valueOf();

Argument of type 'null' is not assignable to parameter of type 'number'. ts(2345)

View Problem No quick fixes available

console.log(getValue(null));



```
Types in TypeScript
```

```
function getValue(x: number | null) {
    return x.valueOf();
```

Object is possibly 'null'. ts(2531)

(parameter) x: number | null

View Problem No quick fixes available

console.log(getValue(null));

17-214/514



Step Back

• Why do we care about types so much?



Step Back

- Why do we care about types so much?
 - We care about *common mistakes*
 - Type errors happen to be very common
 - What else is common?



Step Back

- Why do we care about types so much?
 - We care about *common mistakes*
 - Type errors happen to be very common
 - What else is common?
 - Nullity errors
 - Missing imports

```
public void forward(String sender) {
```

```
if (sender <mark>==</mark> "me") {
```

```
sendSelf();
```

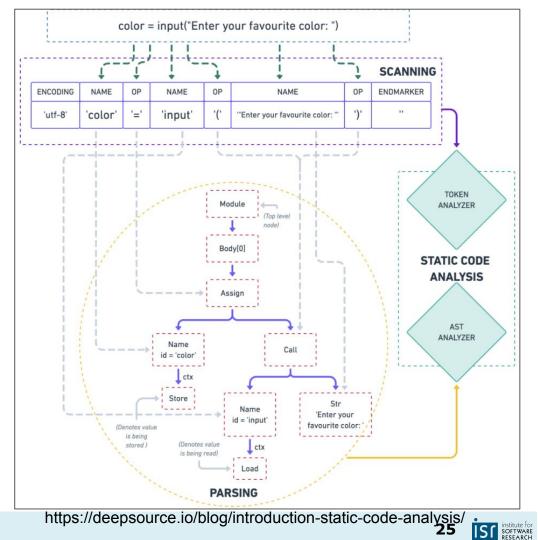
```
} else <mark>if</mark> (sender <mark>==</mark> "other") {
```



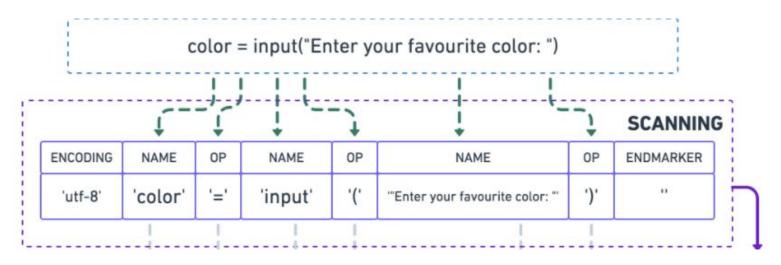
- Detect real or plausible bugs based on code patterns
 - Plausible: look for risk-prone areas
 - Deeply nested loops
 - Overly general types (e.g,. 'any' in TS)
 - Dead code/unused variables
 - Any other places we often make mistakes?



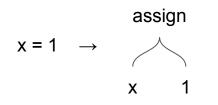
- How?
 - Program analysis +
 Vocabulary of patterns



- Step 1: Tokenization
 - Tokens are like the words of software
 - Lexical categories, incl. punctuation, identifiers, operators, strings

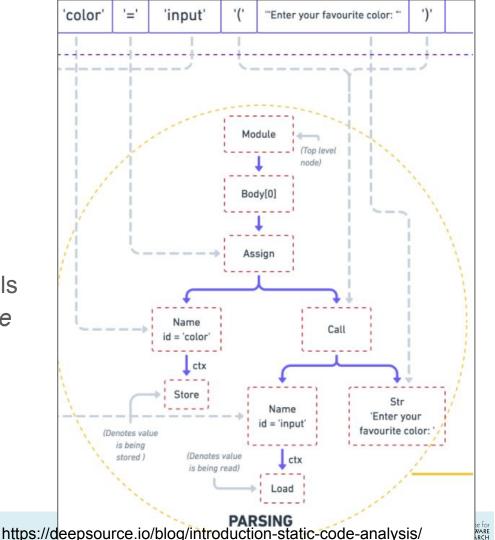


- Step 2: Parsing
 - To the compiler/interpreter, software is a <u>tree</u>
 - Root node is file/module
 - Leaves mainly identifiers, literals
 - Internal nodes capture structure

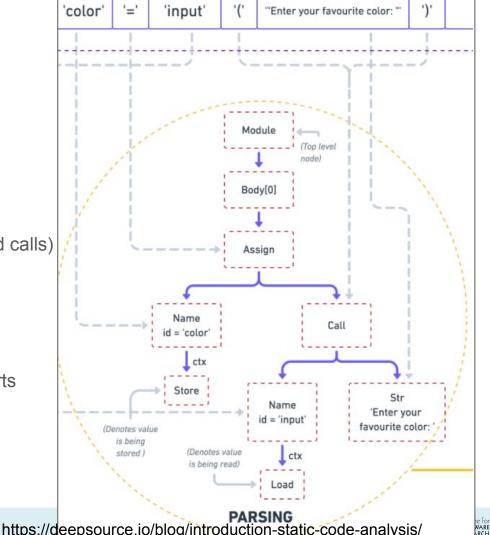


Consider checking out: https://ast.carlosroso.com/

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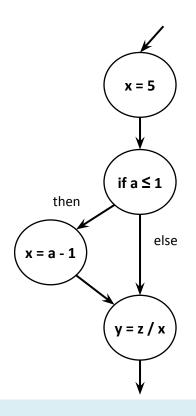
- Step 2: Parsing
 - What does this get us?
 - Rich structure
 - Syntactic types (variables, method calls)
 - Dead code, deep nesting
 - A lot of type resolution
 - What vars are stored, loaded
 - Not complete!
 - Need to *build* to understand imports



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- Step 2b: Advanced Analysis
 - The compiler doesn't stop at parsing
 - Familiar?

```
public boolean div(int a, int z) {
    int x = 5;
    if (a <= 1) {
        x = a - 1;
    }
    return z / x;
}</pre>
```





- Step 2b: Advanced Analysis
 - The compiler doesn't stop at parsing
 - There is <u>a lot</u> more down this rabbit hole
 - Control/data-flow, abstract interpretation, (dynamic) symbolic execution,
 - Consider a programming languages, compilers, or program analysis course



- Step 3: register analyzers
 - At the core: walk the tree

```
class ListDefinitionChecker(BaseChecker):
    msg = "usage of 'list()' detected, use '[]' instead"
```

```
def visit_Call(self, node):
```

name = getattr(node.func, "id", None)

if name and name == list.__name__ and not node.args:

self.violations.append((self.filename, node.lineno, self.msg))



- Step 3: register analyzers
- Classic: walk a tree \rightarrow
- Modern: build a database of code facts, express analysis as queries over that database.
 - This is how CodeQL works!

```
class UnusedImportChecker(BaseChecker):
    def __init__(self):
        self.import_map = defaultdict(set)
        self.name_map = defaultdict(set)
```

```
def _add_imports(self, node):
    for import_name in node.names:
        # Store only top-level module name ("os.path" -> "os").
        # We can't easily detect when "os.path" is used.
        name = import_name.name.partition(".")[0]
        self.import_map[self.filename].add((name, node.lineno))
```

```
def visit_Import(self, node):
    self._add_imports(node)
```

```
def visit_ImportFrom(self, node):
    self._add_imports(node)
```

```
def visit_Name(self, node):
    # We only add those nodes for which a value is being read from.
    if isinstance(node.ctx, ast.Load):
        self.name_map[self.filename].add(node.id)
```



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- Modern: build a database of code facts, express analysis as queries over that database.
- This is how CodeQL works!

0]JavaConverter.java			
	t 0b <u>-</u>	iec'	<pre>static Object deserialize (InputStream is) ows IOException { tInputStream ois = new ObjectInputStream(is); n ois.readObject();</pre>	
D	UnsafeDeserialization.ql			
	<pre>from DataFlow::PathNode source, DataFlow::PathNode sink, UnsafeDeserializationConfig conf where conf.hasFlowPath(source, sink) select sink.getNode().(UnsafeDeserializationSink) .getMethodAccess(), source, sink, "Unsafe deserialization of \$@.", source.getNode(), "user input" QL Query Results</pre>			
	alerts ~			
		Ur	Unsafe deserialization of <u>user input</u> .	
	\sim	Path		
		1	<pre>getContent() : InputStream</pre>	
		2	<pre>getContentAsStream() : InputStream</pre>	
		3	<pre>toBufferedInputStream() : InputStream</pre>	
		4	<pre>getInputStream() : InputStream</pre>	
		5	<u>is : InputStream</u>	
		6	ois	
	>	Pa	th	
	> :=	Ur	safe deserialization of user input.	

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- Compared to Linters:
 - Linters mainly enforce style -- comments, quotes, idioms
 - This also requires static analysis! Just nothing particularly fancy
 - Some overlap; good conventions help avoid bugs



- Compared to Parsers:
 - Parsers check for syntactic correctness
 - Can catch bugs as well, e.g. missing ";"
 - Parsing is often a key step in static analysis
 - Hard to do right with just text/regexes.
 - Parsing is a platform for further analyses
 - control-flow, data-flow





So... Static Analysis for Everything?

- Can we find every bug?
 - No! Rice's Theorem

"Any nontrivial property about the language recognized by a Turing machine is undecidable." -- Henry Gordon Rice, 1953

 Every static analysis is necessarily incomplete or unsound or undecidable (or multiple of these)



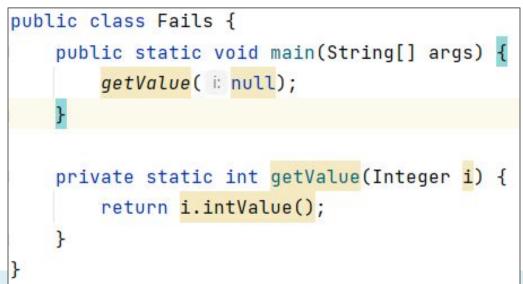
So... Static Analysis for Everything?

- Can we find every bug?
- Can we guarantee correctness?



So... Static Analysis for Everything?

- Can we find every bug?
- Can we guarantee correctness?
 - Yes, but... much less useful

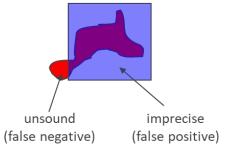


Soundness & Precision

- Since we can't perfectly analyze behavior statically
 - We may miss things by being cautious (unsound; false negative)
 - We might identify non-problems (imprecision, false positive)



Program state covered in actual execution





Program state covered by abstract execution with analysis





The Social Side

• How to deploy tools that are neither sound nor complete?



- Centered around FindBugs (succeeded by SpotBugs)
 - Essentially, a huge collection of risky patterns on Java bytecode
 - Annotated with five levels of concern



- Three experiments in the early 2000s:
 - **1. A dashboard:** run FindBugs overnight, report results in a centralized location

Failed because: dashboard is outside the developer's workflow





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 - 1. A dashboard: run FindBugs overnight, report results in a centralized location

Failed because: dashboard is outside the developer's workflow

 Recurring Fixlt events: company-wide one-week effort to fix warnings Failed because: actually fixed some bugs, but FindBugs is too imprecise (44% of issues were "bugs", but only 16% mattered)



None of these worked!

- Three experiments in the early 2000s:
 - 1. A dashboard: run FindBugs overnight, report results in a centralized location

Failed because: dashboard is outside the developer's workflow

- Recurring FixIt events: company-wide one-week effort to fix warnings Failed because: actually fixed some bugs, but FindBugs is too imprecise (44% of issues were "bugs", but only 16% mattered)
- 3. Add to Code Review: run on every change, allow toggling warnings *Failed because:* too imprecise; suppressing FPs made it inconsistent



Okay so then what?

• What went wrong / what do we need?



Okay so then what?

- What went wrong / what do we need?
 - 1. Precision is key -- developers lose faith in inaccurate tools
 - 2. Provide timely warnings -- in-IDE or rapidly on builds
 - a. Checkers are way more useful during coding
 - 3. Make a platform -- allow adding useful checks



Specifically:

- At compile-time:
 - Perfectly Precise
 - **No** false-positives; never halt a build incorrectly
 - Simple
 - Actionable
 - Ideally to the point of auto-fix suggestions





Specifically:

- At review time: TriCoder
 - 90%+ precise
 - If it drops below, checker gets disabled! Onus on checker authors to fix
 - Actionable, but may require some work
 - Improve correctness or code quality
 - Some compile-time checks moved to review-time!
- Ran 50K times per day -- in 2018



TriCoder

Not
icanalysis;
<pre>ing(), "foo".toString());</pre>

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49

- The gist: Many simple precise checks
 - What else could one do?

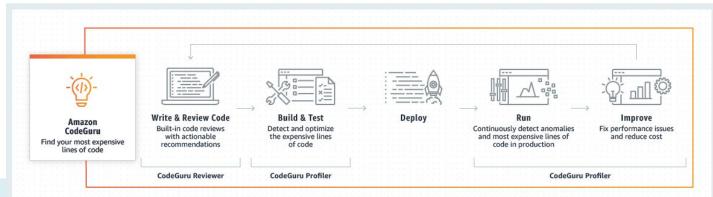


- The gist: Many simple precise checks
 - What else could one do?
- Infer at Facebook
 - Built around separation logic; geared heavily towards tracking resources
 - Null-pointer dereferences, resource leaks, unintended data access
 - Google claims this won't (easily) scale to their multi-billion line mono-repo





- The gist: Many simple precise checks
 - What else could one do?
- Use AI?
 - Rule-mining from previous reviews
 - Detects typical vulnerabilities, bad patterns
 - Mostly fairly simple ML (details limited)



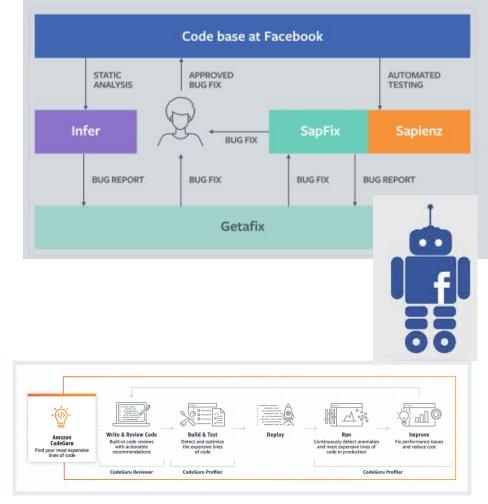
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- The gist: Many simple precise checks
 - What else could one do?
- Use AI?
 - Microsoft's IntelliSense in VSCode
 - Mostly refactorings, code completions
 - Trained on large volumes of code



What else could we do?

- Use more complicated logic One example: Infer, at Facebook (Google claims this won't (easily) scale to their mono-repo.)
- Use AI?
 - Facebook: Getafix, also integrates with SapFix
 - Amazon: CodeGuru
 - Microsoft: IntelliSense in VSCode, mostly refactoring/code completion, trained on large volumes of code
 - Mostly fairly simple ML (details limited)





Summary

- We all constantly make mistakes
 - Static analysis captures common issues
 - Choose suitable abstractions; consider trade-offs
 - E.g., dynamic vs. static typing; sound vs. precise
- At big-tech-scale, automated checks are key
 - Help normalize coding standards
 - Even rare bugs are common at scale
 - But: social factors are very important

